

Bisimilar control systems



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Modeling and Control of Hybrid Systems

Veldhoven, The Netherlands

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http://www.et.tudelft.nl/~disc_ha/



Outline of this mini-course

Lecture 1 : Monday, June 23

Examples of hybrid systems, modeling formalisms

Lecture 2 : Monday, June 23

Transitions systems, temporal logic, refinement notions

Lecture 3 : Tuesday, June 24

Discrete abstractions of hybrid systems for verification

Lecture 4 : Tuesday, June 24

Discrete abstractions of continuous systems for control

Lecture 5 : Thursday, June 26

Bisimilar control systems



Continuous to continuous (Lecture 5)

Lecture 4

T / \approx

$T \equiv T / \approx$

$T \quad y_{t+1} = Fy_t + Gv_t$

Exponential number states

Lecture 5

T / \approx

$T \equiv T / \approx$

$T \quad x_{t+1} = Ax_t + Bu_t$

Bisimilar control systems



Goal

Abstracted Model

$z'_1 = v$
 $z'_2 = z_2$

Observations

$\langle z \rangle = \begin{cases} o_1 & \text{if } z_1 > 5 \wedge z_2 > 5 \wedge z_1 + z_2 < 10 \\ o_2 & \text{if } z_1 \leq -5 \wedge z_2 \leq -5 \wedge z_1 + z_2 > -10 \\ o_3 & \text{otherwise} \end{cases}$

$z_1 = x_1$
 $z_2 = x_4$

Original Model

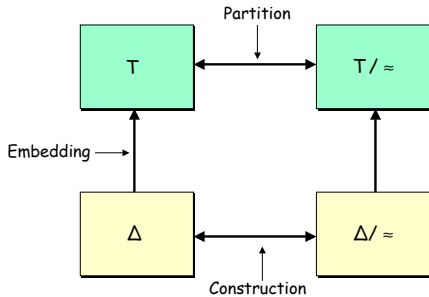
$x'_1 = x_3$
 $x'_2 = -x_2 + x_4 + u$
 $x'_3 = x_1 + x_2$
 $x'_4 = x_4$

Observations

$\langle x \rangle = \begin{cases} o_1 & \text{if } x_1 > 5 \wedge x_4 > 5 \wedge x_1 + x_4 < 10 \\ o_2 & \text{if } x_1 \leq -5 \wedge x_4 \leq -5 \wedge x_1 + x_4 > -10 \\ o_3 & \text{otherwise} \end{cases}$



Game plan



Lossless Embedding

$T_z^\Delta = (Q, \Sigma, \rightarrow, O, \langle \rangle)$

$\Delta \quad x_{k+1} = Ax_k + Bu_k$
 $y_k = Cx_k$

Transition System T_u^Δ

State set $Q = X = \mathbb{R}^n$

Label set $\Sigma = U = \mathbb{R}^m$

Observation set $O = Y = \mathbb{R}^p$

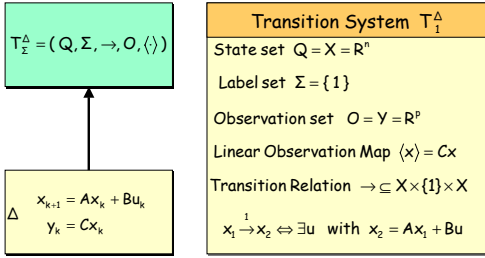
Linear Observation Map $\langle x \rangle = Cx$

Transition Relation $\rightarrow \subseteq X \times U \times X$

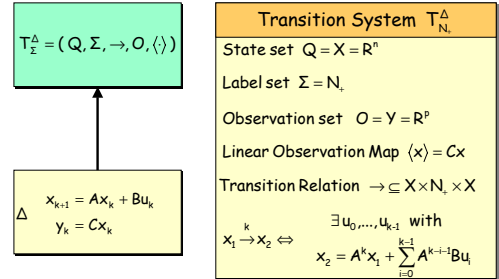
$x_1 \xrightarrow{u} x_2 \Leftrightarrow x_2 = Ax_1 + Bu$



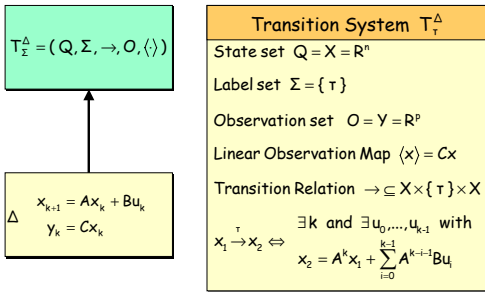
Loose Control...



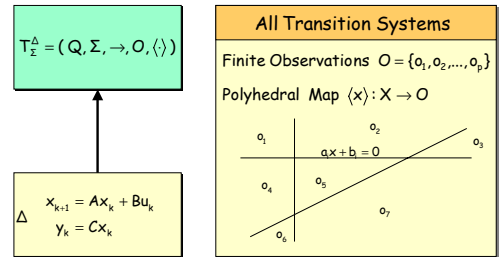
Keep time....



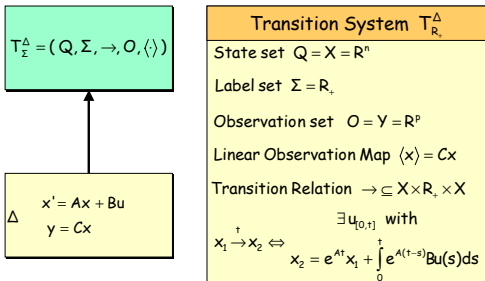
Loose control and time...



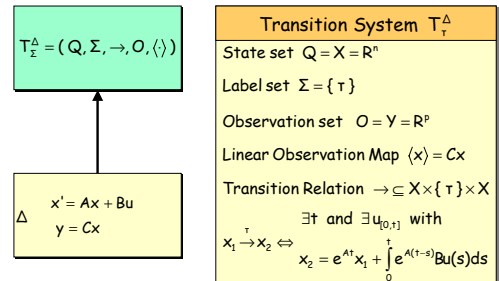
Finite Observations



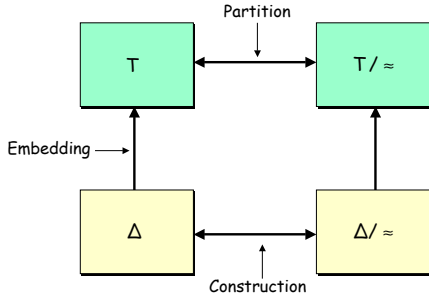
Keep continuous time....



Loose continuous time...



Partitions



Respecting the observations

Two states are equivalent iff

$$x_1 \approx x_2 \Leftrightarrow Hx_1 = Hx_2 \Leftrightarrow x_1 - x_2 \in \text{Ker}(H)$$

for some surjective map $z = Hx$. Simulation $S = (x, Hx)$

Partition is observation preserving iff

Linear observations :

$$\text{Ker}(H) \subseteq \text{Ker}(C)$$

Finite, polyhedral observations :

$$\text{Ker}(H) \subseteq \text{Ker}(a)$$

Respecting the controlled transitions

Respecting the transitions depends on the embedding.

Consider the transition system T_U^Δ

$$\begin{array}{ccc} x_1 & \xrightarrow{u} & x_1' = Ax_1 + Bu \\ \approx & & \approx \\ x_2 & \xrightarrow{u} & x_2' = Ax_2 + Bu \end{array}$$

Proposition : Partition respects the transitions iff
 $A\text{Ker}(H) \subseteq \text{Ker}(H)$

Respecting the timed transitions

Consider the control-abstract transition system T_1^Δ

$$\begin{array}{ccc} x_1 & \xrightarrow{1} & x_1' = Ax_1 + Bu \text{ for some } u \\ \approx & & \approx \\ x_2 & \xrightarrow{1} & x_2' = Ax_2 + Bu' \text{ for some } u' \end{array}$$

Proposition* : Partition respects the transitions iff
 $A\text{Ker}(H) \subseteq \text{Ker}(H) + R(B)$

Similarly

Consider the control-abstract transition system $T_{N_i}^\Delta$

$$\begin{array}{ccc} x_1 & \xrightarrow{k} & x_1' \\ \approx & & \approx \\ x_2 & \xrightarrow{k} & x_2' \end{array}$$

Proposition* : Partition respects the transitions iff
 $A\text{Ker}(H) \subseteq \text{Ker}(H) + R(B)$

Respecting the untimed transitions

Consider the time-abstract transition system T_τ^Δ

$$\begin{array}{ccc} x_1 & \xrightarrow{T} & x_1' \\ \approx & & \approx \\ x_2 & \xrightarrow{T} & x_2' \end{array}$$

Proposition* : Partition respects the transitions iff
 $A\text{Ker}(H) \subseteq \text{Ker}(H) + R(A, B)$

Timed, continuous transitions

Consider the time-abstract transition system $T_{R_t}^\Delta$

$$\begin{array}{ccc} x_1 & \xrightarrow{\dagger} & x_1' \\ \approx & & \approx \\ x_2 & \xrightarrow{\dagger} & x_2' \end{array}$$

Proposition* : Partition respects the transitions iff

$$AKer(H) \subseteq Ker(H) + R(A, B)$$



Untimed, continuous transitions

Consider the time-abstract transition system T_T^Δ

$$\begin{array}{ccc} x_1 & \xrightarrow{T} & x_1' \\ \approx & & \approx \\ x_2 & \xrightarrow{T} & x_2' \end{array}$$

Proposition* : Partition respects the transitions iff

$$AKer(H) \subseteq Ker(H) + R(A, B)$$



Summary

In addition to preserving the observations...

Embedding	Condition
$T_1^\Delta \quad T_{N_t}^\Delta$	$AKer(H) \subseteq Ker(H) + R(B)$
T_T^Δ	$AKer(H) \subseteq Ker(H) + R(A, B)$
$T_{R_t}^\Delta$	$AKer(H) \subseteq Ker(H) + R(A, B)$
T_T^Δ	$AKer(H) \subseteq Ker(H) + R(A, B)$



Coarsest Bisimulation

Find map $z=Hx$ which abstracts as much as possible.
Thus $Ker(H)$ must be maximal but also...

Preserves observations

$$Ker(H) \subseteq Ker(C)$$

Preserves transitions of T_1^Δ

$$AKer(H) \subseteq Ker(H) + R(B)$$

Other variations for other embeddings...



Coarsest Bisimulation Algorithm

Maximal controlled invariant subspace computation

$$V_0 = Ker(C)$$

$$V_{k+1} = V_{k-1} \cap A^{-1}(V_{k-1} + R(B))$$

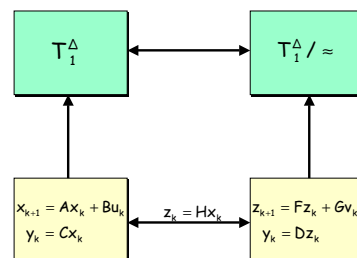
Then $V^* = V_n$ is the maximal desired subspace

Once V^* is computed, then pick map $z=Hx$ such that

$$Ker(H)=V^*$$

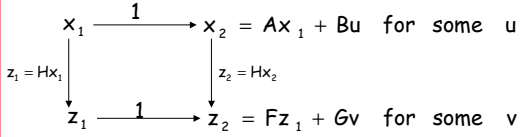


Constructing the abstraction



Construction

Construction of the generator of system T_1^Δ / \approx



Equivalently, for any x, u , there must exist a v such that

$$HAX + HBU = FHx + Gv$$



H-related control systems

Consider discrete-time or continuous-time linear systems

$$(X) \quad x' = Ax + Bu$$

$$(Z) \quad z' = Fz + Gv$$

where $z=Hx$ is surjective. Then (Z) is H-related to (X) if for all x, u there exists v such that

$$H(Ax+Bu) = FHx + Gv$$

Proposition* : Given $x'=Ax+Bu$ and onto map $y=Hx$ choose

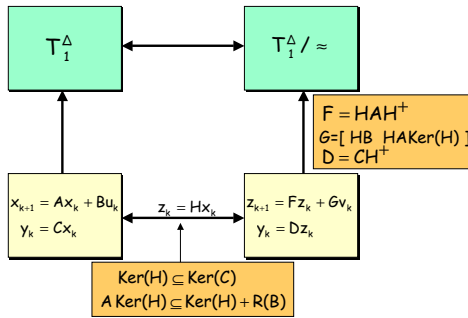
$$F = HAH^+$$

$$G = [HB \quad HAKer(H)]$$

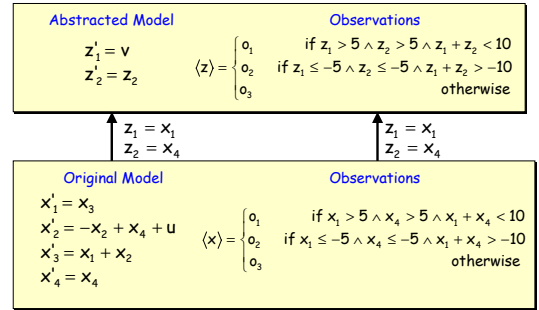
Then $z'=Fz+Gv$ is H-related to $x'=Ax+Bu$



Bisimilar Linear Systems



Goal achieved



Preserving Controllability

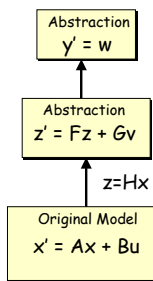
Theorem* : Given onto map $z=Hx$ and linear system $x'=Ax+Bu$ (X) construct the canonical H-related system $z'=Fz+Gv$ (Z) with $F = HAH^+$ and $G = [HB \quad HAKer(H)]$. Assume that $Ker(H) \subseteq R(A, B)$. Then (X) is controllable if and only if (Z) is controllable.

Reachable sets satisfy $R(A, B) = H^{-1}(R(F, G))$

This leads to a hierarchical controllability algorithm...



Hierarchical Controllability Algorithm



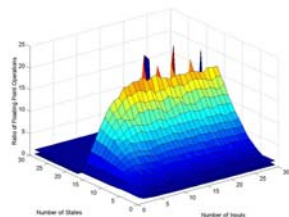
Hierarchical Algorithm

- 0 Initially $x'=Ax+Bu, 0 \leq k < n$
- 1 If rank(B) is
 - 0 : System uncontrollable
 - n : System controllable
- 2 Find H with $Ker(H) = Im[B \ AB \ \dots \ A^{k-1}B]$
- 3 Compute
 - $A := HAH^+$
 - $B := [HB \ HAKer(H)]$
- 4 Go to 2



Hierarchical Controllability Algorithm

Comparison with Kalman rank test

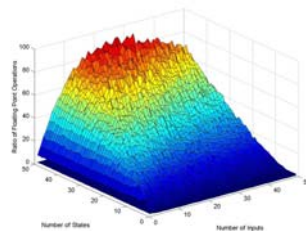


$$\text{Ratio} = \frac{\text{flops of other tests}}{\text{flops of our algorithm}}$$



Hierarchical Controllability Algorithm

Comparison with Popov-Belevitch-Hautus (PBH) test

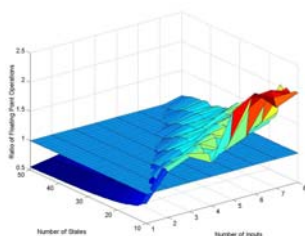


Algorithm with $k=0$ recovers best known algorithm



Hierarchical Controllability Algorithm

Comparison of our Algorithm with $k=0$ and $k=1$



Higher Lie brackets help for underactuated systems



Property Preserving Abstractions

Control is critical for abstraction!

Property	Condition
Controllability	$\text{Ker}(H) \subseteq \text{R}(A, B)$
Output Controllability	$A \text{Ker}(H) \subseteq \text{Ker}(H) + \text{R}(A, B)$
Trajectories	$\text{Ker}(H) \subseteq \text{R}(B)$
Stabilizability	$\text{Ker}(H) \subseteq X^* + \text{R}(A, B)$



Some take home messages

New (hybrid) models, but also new (hybrid) questions

- Partial synchronization of continuous systems
- Logic is entering our world
 - Temporal logic for complicated specifications
 - First-order logic for syntactically specifying hybrid systems

Algorithmic approaches to analysis and controller design

- Is your design method computationally feasible?
- Is your design method computationally efficient?
- Focus on tool development

Decidability boundary for hybrid problems is mature

- Complexity boundary is not

(Bi)simulation relations are very useful

- Theoretically : As a system theoretic concept
- Practically : As a complexity reduction mechanism



Some future directions

Stochastic hybrid systems

Equivalence (model reduction) of hybrid systems

Approximate but efficient algorithms for analysis and design

Understanding compositionality and concurrency

Hybrid (heterogeneous) systems in a broader context

A unified systems theory



An invitation

Hybrid Systems : Computation and Control
University of Pennsylvania
March 25-27, 2004



<http://www.seas.upenn.edu/hybrid/HSCC04/>

