# Distributed Randomized Optimization for Large Scale Interconnected Systems

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April 12, 2017



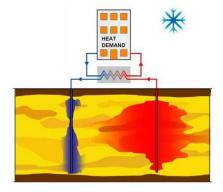


# Aquifer Thermal Energy Storage (ATES)

- A large-scale natural subsurface storage for thermal energy
- An innovative method for thermal energy balance in smart grids

#### Cold season:

- The building requests thermal energy for the heating purpose
- Water is injected into cold well and is taken from warm well
- The stored water contains cold thermal energy for next season



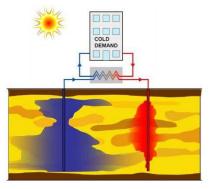
[Rostampour et al., JEP, 2016]

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#### Warm season:

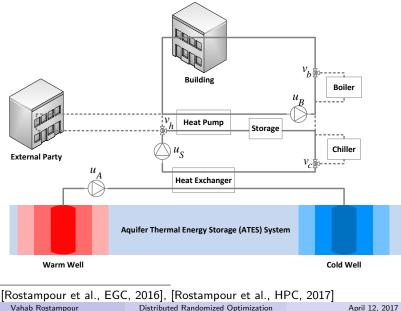
- The building requests thermal energy for the cooling purpose
- Water is injected into warm well and is taken from cold well
- The stored water contains warm thermal energy for next season



[Rostampour et al., JEP, 2016]

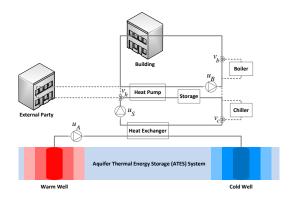
## How to Deal with ATES Systems in Smart Thermal Grids?

# Single Building with ATES System



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# Single Building with ATES System



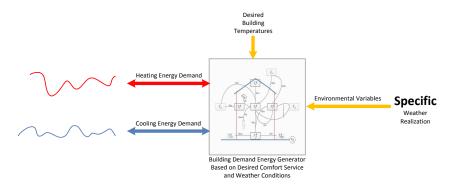
- A complete model of building thermal system integrated with ATES system
- Simulation results for more than a month

[Rostampour et al., EGC, 2016], [Rostampour et al., HPC, 2017]

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- Nonconvex problem formulation
- Computationally intractable optimization problem for long simulation study
- Infeasible setup for the network of ATES systems

# Single Building Thermal Energy Demand

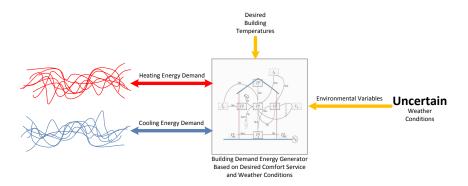


#### **Thermal Energy Demand Profile:**

- Complete and detailed building dynamical model
- Desired building temperatures (local controller unit)
- In specific weather realization, certain demand profiles are generated

[Rostampour & Keviczky, ECC, 2016]

# Single Building Thermal Energy Demand

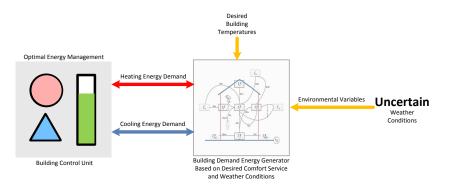


#### **Thermal Energy Demand Profile:**

- Complete and detailed building dynamical model
- Desired building temperatures (local controller unit)
- In uncertain conditions, uncertain demand profiles are generated

[Rostampour & Keviczky, ECC, 2016]

# Single Building Climate Comfort System

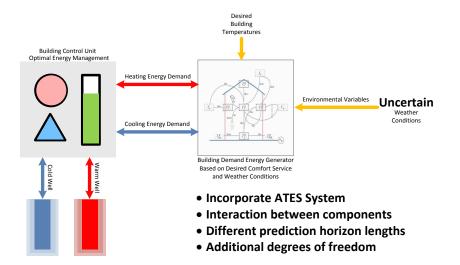


#### **Building Control Unit:**

- Main components: Boiler, HP, HE, micro-CHP, Storage Tank
- ON/OFF status together with production schedule as decisions
- Control Objective: thermal energy balance for the overall systems

[Rostampour & Keviczky, ECC, 2016]

# Building Climate Comfort with ATES Systems



[Rostampour & Keviczky, IFAC, 2017]

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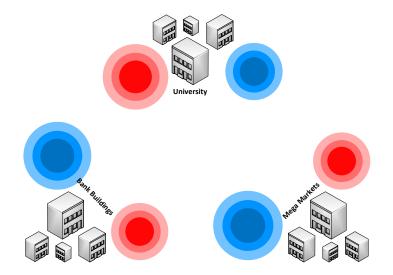






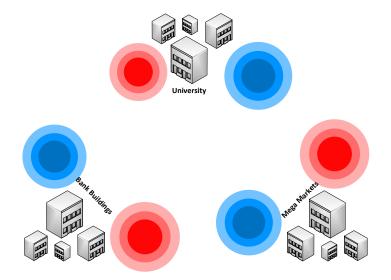
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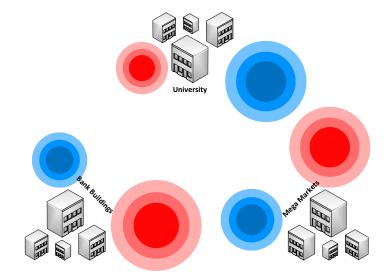
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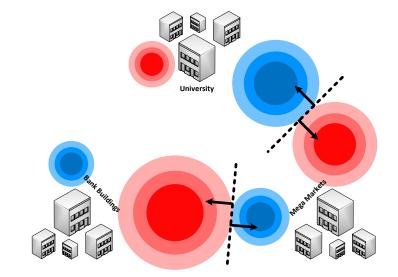
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#### [Rostampour & Keviczky, IFAC, 2017]

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## Achievements & Developments

#### [Rostampour & Keviczky, IFAC, 2017]

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Challenge: Optimizing the performance of a network ...

- **1** Computation: Problem size is too large!
- **2** Communication: Communication bandwidth limitation
- **3** Information Privacy: Agents may not want to share information
- 4 Stochastic Nature:
  - Agents private uncertainty source (local); uncertain thermal energy demand of a single building climate comfort
  - Agents common uncertainty source (shared); uncertain common resource pool, e.g., ATES systems

# Why Distributed?

### Scalable Methodology

- Communication: Only between neighbors
- Computation: Only local; in parallel for all agents

### 2 Preserving Privacy

• Agents do not reveal information about their preferences (encoded by objective and constraint functions) to each other

### 8 Numerous Applications

- Wireless Networks
- Electric Vehicle Charging Control
- Optimal Power Flow with Reserve Scheduling \*
- Energy Management in STGs with ATES Systems <sup>†</sup>

\*[Rostampour et. al., 2017] <sup>†</sup>[Rostampour & Keviczky, 2017

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# Outline

- 1 Centralized Framework
- **2** Distributed Framework
- **3** Soft Communication Scheme
- 4 Case Study: STGs with ATES Systems
- **5** Conclusions and Future work

### Centralized Deterministic Program

$$\begin{array}{rcl} \min_{x} & \sum_{i} f_{i}(x) & \longrightarrow & f_{i}(\cdot) & : & \text{objective function of agent } i \\ \text{s.t.} & x \in \mathcal{X}_{i} \ , \ \text{for all } i & \longrightarrow & \mathcal{X}_{i} & : & \text{constraint set of agent } i \end{array}$$

### Centralized Deterministic Program

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• How one can deal with uncertain  $\mathcal{X}_i(\delta)$ ?

### Centralized Robust Program

$$\begin{split} \min_{x} & \sum_{i} f_{i}(x) \\ \text{s.t.} & x \in \bigcap_{i} \mathcal{X}_{i}(\delta) \ , \ \text{for all} \ \delta \in \Delta \end{split}$$

### Centralized Robust Program

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Stochastic Setting:

- $\delta$  : Uncertain parameter  $\delta \sim \mathbb{P}$
- $\Delta$  : Possibly unknown distribution and unbounded set
- Semi-infinite optimization problem

#### Centralized Robust Program

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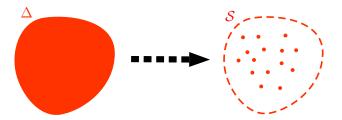
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Centralized Scenario Program

$$\min_{x} \quad \sum_{i} f_{i}(x)$$
s.t.  $x \in \bigcap_{i} \bigcap_{\delta \in \Delta} \mathcal{X}_{i}(\delta)$ 

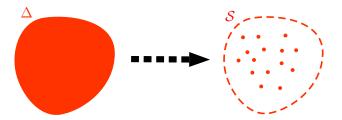
Replace  $\Delta$  with S:



Centralized Scenario Program

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Replace  $\Delta$  with S:



Probabilistic Feasibility Certificate

Centralized Scenario Program  $\mathcal{P}_{\mathcal{S}}$  Centralized Stochastic Program  $\mathcal{P}_{\Delta}$ 

$$\begin{split} \min_{x} & \sum_{i} f_{i}(x) & \min_{x} & \sum_{i} f_{i}(x) \\ \text{s.t.} & x \in \bigcap_{i} \bigcap_{\delta \in \mathcal{S}} \mathcal{X}_{i}(\delta) & \text{s.t.} & \mathbb{P}\Big(\delta \in \Delta \, : \, x \notin \bigcap_{i} \mathcal{X}_{i}(\delta)\Big) \leq \varepsilon \end{split}$$

• Is  $x^*_{\mathcal{S}} \vDash \mathcal{P}_{\mathcal{S}}$  feasible for  $\mathcal{P}_{\Delta}$ ?

• Is this true for any S?



Probabilistic Feasibility Certificate

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Probabilistic Feasibility Certificate

Centralized Scenario Program 
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\end{array}$$

Probabilistic Feasibility [Calafiore & Campi, TAC 2006] Fix  $\beta \in (0, 1)$  and S, then

$$\mathbb{P}^{|\mathcal{S}|}\left(\mathcal{S}\in\Delta^{|\mathcal{S}|} : \mathbb{P}\left(\delta\in\Delta : x^*_{\mathcal{S}}\notin\bigcap_i\mathcal{X}_i(\delta)\right)\leq\varepsilon(d,|\mathcal{S}|,\beta)\right)\geq 1-\beta$$

Probabilistic Feasibility Certificate

Probabilistic Feasibility

Fix  $\beta \in (0,1)$  and  $\mathcal{S}$ , then

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Complexity of  $\varepsilon(d, |\mathcal{S}|, \beta)$ :

- Logarithmic in  $\beta$ :  $\beta$  can be set close to "zero"
- Linear in  $|\mathcal{S}|^{-1}$ : the more data the better the result
- Linear in *d*: number of samples from *S* which "support" the solution, i.e. would leave it unchanged (# decision variables)

# Outline

- Centralized Framework
- **2** Distributed Framework
- **3** Soft Communication Scheme
- **4** Case Study: STGs with ATES Systems
- **5** Conclusions and Future work

There are two possible cases:

• Private (local) uncertainty source: i.e. uncertain thermal energy demand of a single building climate comfort

 $x_i \in \mathcal{X}_i(\delta_i)$ , for all  $\delta_i \in \Delta_i$  and for all agents i

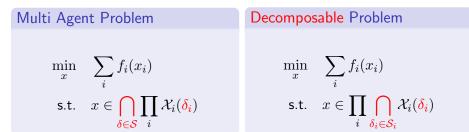
2 Common uncertainty source: i.e. uncertain common resource pool between neighboring agents

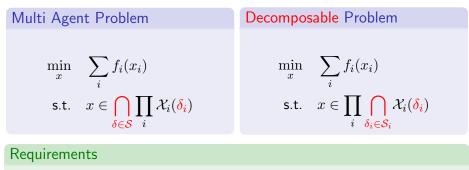
$$x \in \bigcap_k \mathcal{X}_{c_k}(\delta_{c_k}) \;, \; \text{for all } \delta_{c_k} \in \Delta_{c_k}$$

## Private Uncertainty Source

### Multi Agent Problem

$$\min_{x} \quad \sum_{i} f_{i}(x_{i})$$
s.t.  $x \in \bigcap_{\delta \in S} \prod_{i} \mathcal{X}_{i}(\delta_{i})$ 





• Decomposable uncertainty source:

$$\delta := [\delta_1, \cdots, \delta_i, \delta_j, \cdots]$$
 and  $\mathcal{S} := \prod_i \mathcal{S}_i$ 

#### Single Agent Problem

Probabilistic Feasibility for Single Agent Problem<sup>\*</sup> Fix  $\varepsilon_i \in (0, 1)$ ,  $\beta_i \in (0, 1)$  and  $S_i$ , then

$$\mathbb{P}^{|\mathcal{S}_i|}\left(\mathcal{S}_i \in \Delta_i^{|\mathcal{S}_i|} : \mathbb{P}\left(\delta_i \in \Delta_i : x^*_{\mathcal{S}_i} \notin \mathcal{X}_i(\delta_i)\right) \le \varepsilon_i\right) \ge 1 - \beta_i$$

\*[Calafiore & Campi, TAC 2006]

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Distributed Randomized Optimization

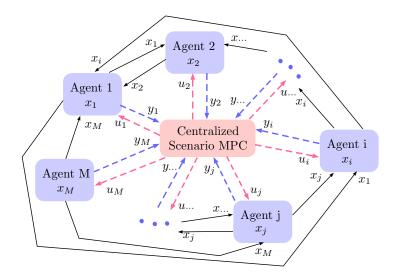
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Probabilistic Feasibility for Multi Agent Problem<sup>\*</sup>  
If 
$$\varepsilon = \sum_{i} \varepsilon_{i} \in (0, 1), \ \beta = \sum_{i} \beta_{i} \in (0, 1) \text{ and given } S$$
, then  
 $\mathbb{P}^{|S|} \left( S \in \Delta^{|S|} : \mathbb{P} \left( \delta \in \Delta : x_{S}^{*} \notin \prod_{i} \mathcal{X}_{i}(\delta_{i}) \right) \leq \varepsilon \right) \geq 1 - \beta$ 

\*[Rostampour & Keviczky, 2017]

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Dynamically Coupled Systems

#### Centralized Scenario Program

$$\begin{split} \min_{\{\boldsymbol{u}_i\}_{\forall i \in \mathcal{N}}} & \sum_{i \in \mathcal{N}} f_i(x_{i,k}, u_{i,k}) \\ \text{s.t.} & x_{i,k+1}^{(i)} = A_{ii} x_{i,k}^{(i)} + B_i u_{i,k} + C_i \delta_{i,k}^{(i)} + \sum_{j \in N_i} A_{ij} x_{j,k}^{(i)} , \ x_{i,k}^{(i)} = x_{i,0} \\ & x_{i,k+\ell}^{(i)} \in \mathcal{X}_i \ , \quad \forall \ell \in \mathbb{N}_+ \ , \quad \forall \delta_{i,k}^{(i)} \in \mathcal{S}_{\delta_i} \\ & u_{i,k} \in \mathcal{U}_i \ , \quad \forall k \in \mathcal{T} \ , \forall i \in \mathcal{N} \end{split}$$

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Dynamically Coupled Systems

Distributed Scenario Program

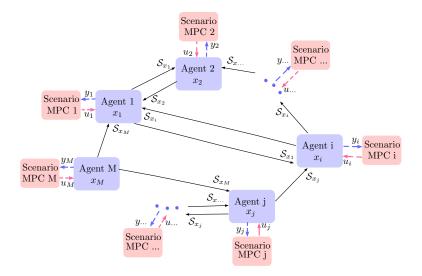
$$\begin{split} \min_{\substack{\{u_{i,k}\}_{\forall k \in \mathcal{T}}}} & f_i(x_{i,k}, u_{i,k}) \\ \text{s.t.} & x_{i,k+1}^{(i)} = A_{ii} x_{i,k}^{(i)} + B_i u_{i,k} + q_{i,k}^{(i)} , \ x_{i,k}^{(i)} = x_{i,0} \\ & x_{i,k+\ell}^{(i)} \in \mathcal{X}_i \ , \quad \forall \ell \in \mathbb{N}_+ \ , \quad \forall q_{i,k}^{(i)} \in \mathcal{S}_{q_i} \\ & u_{i,k} \in \mathcal{U}_i \ , \quad \forall k \in \mathcal{T} \end{split}$$

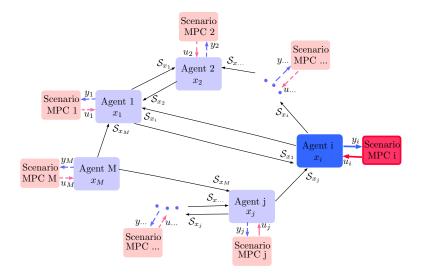
#### Requirements

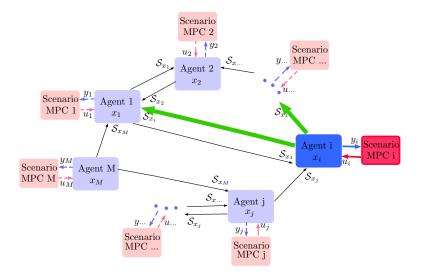
$$\mathcal{S}_{q_i} = \left\{ q_{i,k}^{(i)} : \ q_{i,k}^{(i)} = C_i \delta_{i,k}^{(i)} + \sum_{j \in N_i} A_{ij} x_{j,k}^{(i)} , \ \forall \delta_{i,k}^{(i)} \in \mathcal{S}_{\delta_i} \ , \ \forall x_{j,k}^{(i)} \in \mathcal{S}_{x_j} \right\}$$

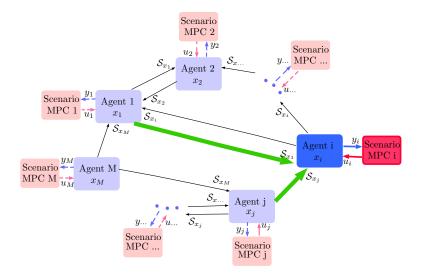
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Distributed Randomized Optimization









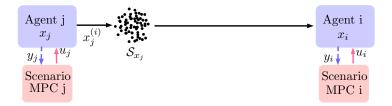
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# Communication Scheme

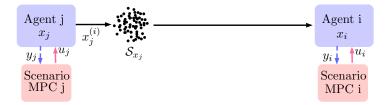
We consider two way to communicate between neighboring agents:

• Hard communication scheme: i.e. agent j has to send exactly the set  $S_{x_i}$  with cardinality  $N_{s_i}$  as it is requested by agent i



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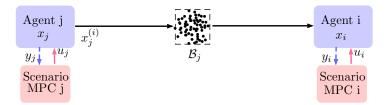
Soft communication scheme: i.e. agent j sends a parametrized set B̃<sub>j</sub> with its desired level of reliability α̃<sub>j</sub> to agent i



▶ It is an interest of agent j to decide about number of scenarios  $N_j$  ◄

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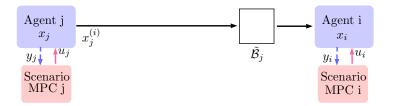
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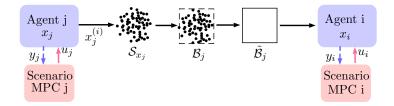
We consider two way to communicate between neighboring agents:

Soft communication scheme: i.e. agent j sends a parametrized set 
 *B˜<sub>j</sub>* with its desired level of reliability 
 *α˜<sub>j</sub>* to agent i



▶ It is an interest of agent j to decide about number of scenarios  $N_j$  ◄

### Soft Communication Scheme



#### Definition

A set  $\tilde{\mathcal{B}}_j \subseteq \mathbb{R}^{m_j}$  is  $\tilde{\alpha}_j$ -reliable if

$$\mathbb{P}\left(\boldsymbol{x}_{j}\in\mathcal{X}_{j}:\;\boldsymbol{x}_{j}\notin\tilde{\mathcal{B}}_{j}
ight)\leq1-\tilde{lpha}_{j}\;,$$

and we refer to  $\tilde{\alpha}_i$  as the level of reliability of the set  $\tilde{\mathcal{B}}_i$ .

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and we refer to  $\tilde{\alpha}_j$  as the level of reliability of the set  $\tilde{\mathcal{B}}_j$ .

#### Theorem<sup>\*</sup>: how one can determine $\tilde{\alpha}_j$ ?

Fix  $\tilde{\beta}_j \in (0,1)$ ,  $\tilde{N}_{s_i}$  and let

$$\tilde{\alpha}_{j} = \left. \begin{smallmatrix} \bar{N}_{s_{i}} - m_{j} \\ \hline \left( \begin{smallmatrix} \bar{\beta}_{j} \\ (\bar{N}_{s_{i}}) \\ m_{j} \end{smallmatrix} \right) \end{smallmatrix} \right.$$

We then have  $\mathbb{P}\left(\boldsymbol{x}_{j} \in \mathcal{X}_{j} : \boldsymbol{x}_{j} \notin \tilde{\mathcal{B}}_{j}\right) \leq 1 - \tilde{\alpha}_{j}$ , with prob.  $1 - \tilde{\beta}_{j}$ .

\*[Rostampour et. al., 2017]

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# Soft Communication Scheme

Probabilistic Feasibility Certificate

The communicated information are reliable with certain level of probability. How one can accommodate such a probabilistically reliable information in the probabilistic feasibility certificate of the local agent?

#### Theorem\*

Given  $\tilde{\alpha}_j \in (0,1)$  and a fixed  $\alpha_i \in (0,1)$ , the state trajectory of a generic agent *i* is probabilistically  $\bar{\alpha}_i$ -feasible for all  $\delta_i \in \Delta_i$ , i.e.,

$$\mathbb{P}\left(x_{i,k+\ell} \in \mathcal{X}_i, \, \ell \in \mathbb{N}_+\right) \geq \bar{\alpha}_i \,,$$

where  $\bar{\alpha}_i = 1 - \frac{1 - \alpha_i}{\tilde{\alpha}_i}$  such that  $\tilde{\alpha}_i = \prod_{j \in \mathcal{N}_i} (\tilde{\alpha}_j)$ .

\*[Rostampour et. al., 2017]

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# Energy Management Problem

#### **Optimization Problem**

- $\min$  thermal energy imbalance error + cost of equipment operation
- s.t. 1) equipment limits
  - 2) imbalance error dynamics
  - 3) ATES system dynamics + local thermal energy balance
  - 4) coupling constraint on the thermal radius between agents

#### [Rostampour & Keviczky, IFAC, 2017]

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Distributed Randomized Optimization

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- s.t. 1) equipment limits
  - 2) imbalance error dynamics
  - 3) ATES system dynamics + local thermal energy balance
  - 4) coupling constraint on the thermal radius between agents

#### Compact Form — x: Decision Variables\*

$$\min_{x_i} \sum_i f_i(x_i) \longrightarrow f_i(\cdot)$$
 : objective function of agent  $i$ 

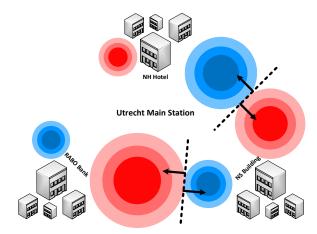
s.t. 
$$x_i \in \mathcal{X}_i(\delta_i)$$
,  $\forall i \longrightarrow \mathcal{X}_i$ : constraint set of agent  $i$   
 $x \in \bigcap \mathcal{X}_{c_k}(\delta_{c_k}) \longrightarrow \mathcal{X}_{c_k}$ : coupling constraint set

\*set-points for control units of buildings and pump flow rate for ATES systems

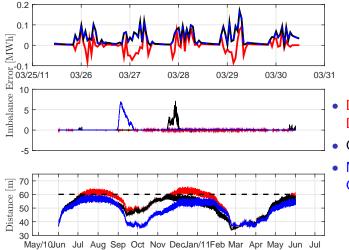
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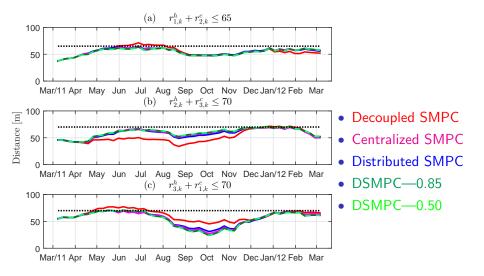
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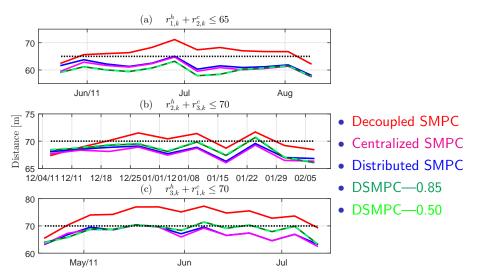


- Three buildings in Utrecht city with real parameters
- Real weather condition data from 2010 to 2012



- Deterministic
   Decoupled MPC
- Centralized SMPC
- Move-Blocking Centralized SMPC





# Outline

- Centralized Framework
- ② Distributed Framework
- **3** Soft Communication Scheme
- Gase Study: STGs with ATES Systems
- **5** Conclusions and Future work

- Distributed randomized optimization to deal with private (local) uncertainty source over a network of dynamically coupled systems
- Soft communication scheme with an extension of probabilistic feasibility guarantee
- Application to energy management of smart thermal grids (STGs) with aquifer thermal energy storage (ATES) system

#### What comes next?

- Distributed randomized optimization for weakly coupled problems together with common uncertainty source
- Preserving privacy of individual agents in a network
- Distributed randomized optimization for hierarchical decision making in uncertain dynamical environment
- More applications distributed (multi-area) reserve scheduling and optimal power flow over AC power networks

# Thank you for your attention! Questions?



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Distributed Randomized Optimization

# Distributed Randomized Optimization for Large Scale Interconnected Systems

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